## LEARN DEPENDENTLY. TVPED PROGRAMMING WITHIDRIS



## - Qpuffinfiesh

- Tiny contrihutor to ldris (18 commits)
- Played with dependent types for 2 years
$>$ Been doing dris for 6 months



## > Small experience with Haskell

## Have an install of ldris (can be tridky)

\$ brew install ghc cabal-install
\$ cabal update
\$ cabal install alex
\$ cabal install idris
nodule Maln
tmport favnifeript
main 180 ()
min = alert "llello morid!"

## Contris mandoty

Henilly

Var _ impr__blelnt $\rightarrow$ (tunctianti
/f Cepyright (e) 2tes Tin wu
// All lights heserved
// See "licresse" for details.
// Basic Juaberigt 3 tibrary - subset useful for RSA encryption.
/f Bits per digit
ver dilts!
/f JavaScrigt engiee analysis
yer cavary = txilesabeetcafe!
var 1 _la $=$ (tcanarylixiftrffl-sbefcafel:
/f (public) Censtractar
fusction Bigintmgeria,b,el
1f(a) I= Nilli!
If ("manber" man trpeof in) this.frankuber(a, b, c)

## OUTLNE

1. Overview of dependent types and Idris 2. Work through exercises, I lead 3. Work through exercises, Ihelp

# Bad neNs: most softwave cannot be reasoned about - Paul Phillips 

## - Curry-Howarid; programs are proofs - Let's make our proofs interesting

- Therefore let's use a powerful type system


## MISGONGEPTIONS

## D Itris is harider than Haskell $\checkmark$ Dependent types are hard

# DEPENDENT TMPES EVERTHING IS A TERW 

isIdris : Bool
isIdris = True
one : Nat
one = if isIdris then S Z else Z

StringList : Type
StringList = if isIdris then List Char else Int

- Types and Kinds are values in universes - Types can depend on values
- Free polymorphism, type constructors
the : ( $\mathrm{t}:$ Type) $->(\mathrm{x}: \mathrm{t})->\mathrm{t}$ the $-a=a$
one : Nat one $=$ the Nat Z
id1 : \{t : Type\} -> (x : t) -> t
id1 \{t\} a = a
id2 : (x : t) -> t
id2 a = a
id3 : t -> t
id3 a = a


## Option : Type -> Type

 Option = Maybe
\$ idris --total
\$ idris --warnpartial
\%default total
total plusOne : Nat -> Nat plusOne Z = S Z
plusOne (S n) = S (S n)

I am often asked' 'how do I inplement a sorper as a program in your terminating language? - Conor McBirite

## I reply that I do not: a

 serler is a coprogram in a language guarartecing lovencess - Conor McBirite
# - We always make progress <br> > Watch out for the totality checker! $>$ Church-Rosser theorem <br> - Evaluation is really normalisation! - Gan still do it all! 


data (=) : a -> b -> Type where refl : x = x
$x: 1=1$
$x$ = refl
$y: 1+1=2$
y = refl

X : \{a : Nat \} -> a - a = Z
X \{a=Z\} = refl
$x$ \{a=S k\} = X \{a=k\}
y : \{a : Nat \} -> a - a = Z
y \{a\} = replace $\{P=1 x=>(a-x=Z)\}$
(plusZeroRightNeutral a)
(minusPlusZero a Z)

X : \{a : Nat \} -> a - a = Z
x = ?xproof
xproof = proof
intros
rewrite (minusPlusZero a Z)
rewrite (plusZeroRightNeutral a) trivial

# - The probilem of dependent types $\checkmark$ Values are unified 

## - Checked for syntactic/terme equality

## WIIY DRIN?

## - LIVM, C, Java, IS hackends

> > FFI

Lots of syntatitic sugar > Tactic rewriting

- Allows more lying/cheating
- REPL, editor modes, doc tools


Version 0.9.9.2
http://www.idris-lang.org/
Type :? for help
instance Show Bits64×2 where
show $\mathrm{X}=\mathrm{a}$
case viewB64×2 x of -
(a, b) $\Rightarrow>$
"<" ++ prim_toStrB64 a
+ " $^{+}$, " + prim_toStrB64 b
$++\quad ">$ "
instance (Show a, Show b) $\Rightarrow$ Show ( $a$, b) where
show $(\mathrm{x}, \mathrm{y})=$ " $("++$ show $\mathrm{x}++$ ", " ++ show $\mathrm{y}++$ ")"
instance Show $a \Rightarrow$ Show (List a) where
show xs = "[" ++ show' "" xs ++ "]" where
show' acc [] =acc
show' acc $[\mathrm{x}] \quad=$ acc ++ show x
show' acc (x : xs) = show' (acc ++ show x ++ ", ") xs
instance Show a $\Rightarrow$ Show (Vect n a) where -r
show xs = "[" ++ show' xs ++ "]" where
show' : Vect n a $->$ String
show' [] = ""
show' $[x]=$ show $x$
show' ( $\mathrm{x}: \mathrm{x}$ x ) = show $\mathrm{x}++$ ", " ++ show' xs
instance Show $\mathrm{a} \Rightarrow$ Show (Maybe a) where -
show Nothing $=$ "Nothing"
show (Just x ) $=$ "Just " ++ show x
---- Functor instances
instance Functor Primio where
map $f$ io $=$ prim_io_bind io (prim_io_return f) -
instance Functor 10 where -
map f io $=$ io_bind io ( b => io_return (f b)) =
instance Functor Maybe where
map $f$ (Just $x$ ) $=$ Just ( $f x$ )
$\operatorname{map} f$ Nothing $=$ Nothing
instance Functor (Either e) where-
map $f$ (Left 1 ) $=$ Left 1
map $f$ (Right $r$ ) $=$ Right ( $\mathbf{f r}$ )
---- Applicative instances
instance Applicative PrimIo where
pure = prim_io_return
Prelude.idr
data Paritu ：Nat－＞Tupe where Even ：〔п ：Nat〕－＞Paritu（п＋п） Odd ：（п ：Nat）－＞Parity〔（S 〔п＋п〕）

## HOW TO DRIS

## - Idris Tutorial

 $>$ Idris Ilibrary docs - Heris Ilibrary source
## Beginning Haskell: a Project Based Approach

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## $\downarrow$ printf

## - Equality proofs

$>$ Verified algehra
$>$ Vector filtering

## http://goo.gl/gfdne

